Resource theories in the asymptotic and catalytic regime

Tobias Fritz

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Asymptotics in resource theories

- \triangleright In resource theories, asymptotics are usually easy.
- ▷ It makes sense to study asymptotics first.
- ▷ Historical examples:
 - ▷ Monge 1781: **optimal transportation**.
 - ▷ Carnot 1824: **Carnot efficiency**.
 - ▷ Kantorovich 1939: linear programming.
 - ▷ Shannon 1948: channel coding theorem.
- ▷ For example, Shannon's theorem is asymptotic and has resulted in the development of coding theory.

Majorization

Let

 $p = (p_1 \ge \ldots \ge p_n > 0), \qquad q = (q_1 \ge \ldots \ge q_m > 0)$

be probability vectors.

Proposition

The following are equivalent:

 \triangleright There is a bistochastic matrix T such that

$$Tp = q.$$

 \triangleright For all k,

$$\sum_{i=1}^k p_i \ge \sum_{i=1}^k q_i.$$

When these hold, we say that p majorizes q,

 $p \succeq q$.

Majorization

 \triangleright Intuition: *q* contains more randomness than *p*.

Theorem (Nielsen '99)

For bipartite pure-state entanglement,

$|\psi\rangle \stackrel{\rm LOCC}{\longrightarrow} |\phi\rangle$

if and only if the spectra of the reduced density matrices display majorization,

 $|\psi\rangle\langle\psi|_{A}^{\downarrow} \succeq |\phi\rangle\langle\phi|_{A}^{\downarrow}.$

- ▷ Provides an easy-to-use criterion for LOCC.
- ▷ What about the asymptotics?
- $\triangleright \text{ Since } |\psi\rangle \mapsto |\psi\rangle \langle \psi|_A^{\downarrow} \text{ preserves tensor products, it's enough to consider asymptotic majorization.}$

The Rényi entropies

$$H_{lpha}(p) \coloneqq rac{1}{1-lpha} \log\left(\sum_{i} p_{i}^{lpha}
ight)$$

will play an important role.

Special cases:

$$H_0(p) = \log |\{i \mid p_i > 0\}|$$
$$H_1(p) = -\sum_i p_i \log p_i$$
$$H_{\infty}(p) = -\log \max_i p_i.$$

Asymptotic majorization

Theorem (Jensen '19) If $H_{\alpha}(p) > H_{\alpha}(q) \quad \forall \alpha \in [0, \infty],$ then $p^{\otimes n} \prec q^{\otimes n} \quad \forall n \gg 1.$

Conversely, $p^{\otimes n} \leq q^{\otimes n}$ for some *n* implies $H_{\alpha}(p) \geq H_{\alpha}(q)$.

▷ Implies a rate formula (TF '17):

$$R(|\psi\rangle \to |\phi\rangle) = \inf_{\alpha} \frac{H_{\alpha}(|\psi\rangle\langle\psi|_{A}^{\downarrow})}{H_{\alpha}(|\phi\rangle\langle\phi|_{A}^{\downarrow})}$$

Catalytic majorization

Theorem (Klimesh '07, Turgut '07) If $H_{\alpha}(p) > H_{\alpha}(q) \quad \forall \alpha \in [0, \infty],$ then $\exists r : p \otimes r \leq q \otimes r.$ Conversely, $p \otimes r \leq q \otimes r$ implies $H_{\alpha}(p) \geq H_{\alpha}(q).$

▷ Why are asymptotic and catalytic majorization essentially equivalent?

 \triangleright And what's special about the H_{α} ?

 $\triangleright\,$ I will answer these questions and explain the general theorem.

▷ The Rényi entropies are additive monotones:

 \triangleright Additivity:

$$H_{\alpha}(p\otimes q)=H_{\alpha}(p)+H_{\alpha}(q).$$

▷ Monotonicity:

$$p \preceq q \implies H_{\alpha}(p) \geq H_{\alpha}(q).$$

 \triangleright There are other additive monotones, such as $H_0 + H_1$.

 \triangleright So something is still missing.

▷ Instead of the H_{α} , consider just

$$\|p\|_{\alpha} = \sum_{\alpha} p_i^{\alpha}$$

for $\alpha \neq 1, \infty$.

 \triangleright We can use $\|p\|_{\alpha} > \|q\|_{\alpha}$ instead of $H_{\alpha}(p) > H_{\alpha}(q)$.

 \triangleright The $\|\cdot\|_{\alpha}$ satisfy multiplicativity

 $\|p\otimes q\|_{lpha} = \|p\|_{lpha} \|q\|_{lpha}$

and monotonicity,

$$p \leq q \implies \|p\|_{\alpha} \geq \|q\|_{\alpha}.$$

D They also satisfy additivity under direct sum

 $\|p \oplus q\|_{lpha} = \|p\|_{lpha} + \|q\|_{lpha}$

if we allow unnormalized probability vectors.

▷ Therefore they are monotone semiring homomorphisms

Major $\longrightarrow \mathbb{R}_+$.

Definition

An ordered semiring $(S, +, \cdot, \geq)$ is an algebraic structure satisfying the usual equations, and

$$x \ge y \implies x+z \ge y+z, \quad xz \ge yz,$$

 \triangleright Major is ordered semiring of probability vectors with $(\oplus, \otimes, \preceq)$.

- ▷ But what about H_1 and H_∞ ?
- \triangleright For H_{∞} , consider instead

$$\|p\|_{\infty}=\max_{i}p_{i}.$$

- ▷ Still a multiplicative monotone.
- \triangleright But now

$$\|p \oplus q\|_{\infty} = \max(\|p\|_{\infty}, \|q\|_{\infty}).$$

Definition

The tropical reals are the ordered semiring

$$\mathbb{TR}_+ \coloneqq ([0,\infty), \max, \cdot, \geq).$$

- \triangleright What about H_1 ?
- ▷ No well-behaved "exponential" exists.
- ▷ Still have additivity

$$H_1(p\oplus q)=H_1(p)+H_1(q)$$

and monotonicity.

▷ Not the usual "additivity" of entropy!

 $\triangleright~$ On tensor products, we have the Leibniz rule

 $H_1(p \otimes q) = H_1(p) ||q||_1 + ||p||_1 H_1(q).$

We say that H_1 is a **derivation** at $\|\cdot\|_1$.

To summarize, the following types of monotones are important:

- $\triangleright \ \ \text{Semiring homomorphisms Major} \to \mathbb{R}_+.$
- $\triangleright \mbox{ Semiring homomorphisms Major} \to \mathbb{TR}_+.$
- $\triangleright~$ The derivations at $\|\cdot\|_1.$

Theorem (With Farooq, Haapasalo, Tomamichel) (a) The $\|\cdot\|_{\alpha}$ for $\alpha \neq \infty$ are all the monotone homs Major $\rightarrow \mathbb{R}_+$. (b) $\|\cdot\|_{\infty}$ is the only monotone hom Major $\rightarrow \mathbb{TR}_+$. (c) H_1 is (essentially) the only derivation at $\|\cdot\|_1$.

hinspace Sketch: Let $f: \mathsf{Major}
ightarrow \mathbb{R}_+$ be a monotone hom. Then

$$f(p_1,\ldots,p_n)=f(p_1)+\cdots+f(p_n).$$

So it's enough to look at vectors of length 1!

> On those, we have the Cauchy functional equation

$$f(pq) = f(p)f(q),$$

whose only well-behaved solutions are the power functions $f(p) = p^{\alpha}$.

General case

- ▷ Now let's generalize to a statement that should apply to many other resource theories too!
- $\triangleright\,$ So let S be any suitably well-behaved ordered semiring and

 $\|\cdot\|\ :\ S\longrightarrow \mathbb{R}_+$

any homomorphism such that

$$p \leq q \implies ||p|| = ||q||.$$

▷ Let us say that the **relevant monotones** are those described above:

- $\triangleright \ \ \mathsf{Monotone} \ \ \mathsf{homs} \ \ \mathcal{S} \to \mathbb{R}_+.$
- $\triangleright \ \text{Monotone homs } S \to \mathbb{TR}_+.$
- $\triangleright \ \ \text{Monotone derivations} \ S \to \mathbb{R} \ \text{at} \ \| \cdot \|.$

Theorem (Vergleichsstellensatz) Let nonzero $x, y \in S$ satisfy ||x|| = ||y||. If f(x) > f(y)

for all relevant monotones f, then:

 \triangleright There is nonzero *c* such that

$$xc \ge yc$$
.

 \triangleright If x is sufficiently generic,

$$x^n \ge y^n \qquad \forall n \gg 1.$$

Conversely, if $xc \ge yc$ for nonzero c or $x^x \ge y^n$ for some $n \ge 1$, then

$$f(x) \ge f(y)$$

for all relevant monotones.

- Provides an almost tight criterion for asymptotic and catalytic convertibility very generally.
- ▷ In particular, shows that asymptotic and catalytic convertibility are essentially equivalent.
- ▷ Recovers known statements on asymptotic and catalytic majorization.
- > Other applications?

Application to representation theory

 $\triangleright\,$ Representations form an ordered semiring with respect to $\oplus, \,\otimes\,$ and containment.

Theorem

For representations of SU(2), the relevant monotones are parametrized by $\alpha \in [0, \infty]$,

$$f_{\alpha}(U) := \sum_{i} \frac{\sinh(lpha \, \dim(U_{i}))}{\sinh lpha}$$

where $U = \bigoplus_{\beta} U_i$ with irreducible U_i ,

Application to representation theory

- ▷ How many spin-¹/₂ qubits do we need in order to simulate n ≫ 1 systems of spin 1, respecting the symmetry?
- \triangleright This can now be computed explicitly: we need Rn, where the rate R is

$$R = R(\operatorname{spin} \frac{1}{2} \to \operatorname{spin} 1) = \inf_{\alpha \ge 0} \frac{\log \frac{\sinh(2\alpha)}{\sinh(a)}}{\log \frac{\sinh(3\alpha)}{\sinh(\alpha)}} = \frac{1}{2}.$$

- ▷ Ordered semirings provide a general framework for resource theories with powerful mathematical results.
- Applying these gives non-constructive proofs for the existence of information processing protocols.
- ▷ No notion of "free resource" or "free operation" is relevant!

Final thoughts 2

Epsilonification is the big open problem: build in approximations in the asymptotics such that e.g. Shannon's channel coding theorem comes out of a general theory as well.

Conjecture

In an epsilonified resource theory, only the derivations are relevant.

▷ E.g. in thermodynamics, this should amount to the free energies, since both entropy and energy are derivations:

$$E(\rho \otimes \eta) = E(\rho) \|\eta\|_1 + \|\rho\|_1 E(\eta).$$

▷ Unfortunately, finding the "right" definitions has turned out to be very difficult.

Quantum thermodynamics^[3]

Consider a system with finite-dimensional Hilbert space \mathbb{C}^d and Hamiltonian H. A state ρ has energy $E(\rho) = \operatorname{tr}(\rho H)$ and entropy $S(\rho) = -\operatorname{tr}(\rho \log \rho)$.

Theorem

For states ρ and σ , the following are equivalent:

(a) There exists an ancilla system of size o(n) with state η and Hamiltonian $H_{\rm anc}$ satisfying $||H_{\rm anc}|| \le o(n)$ and an energy-preserving unitary U with

$$\left\|\operatorname{Tr}_{\operatorname{anc}}\left[U(\rho^{\otimes \mathsf{n}}\otimes\eta)U^{\dagger}\right]-\sigma^{\otimes \mathsf{n}}\right\|_{1}\overset{n\to\infty}{\longrightarrow}0.$$

(b) The states have equal energy and entropy,

$$E(\rho) = E(\sigma), \qquad S(\rho) = S(\sigma).$$

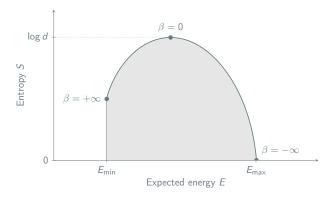
^[1] Carlo Sparaciari, Jonathan Oppenheim, and Tobias Fritz. "A Resource Theory for Work and Heat". In: Phys. Rev. A 96, 052112 (2017). arXiv:1607.01302.

Definition

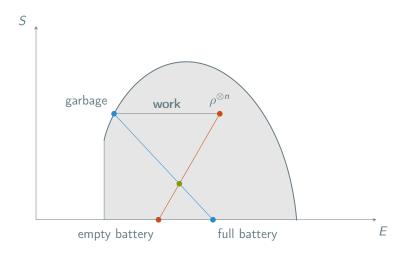
A **macrostate** is an equivalence class of states with respect to asymptotic interconvertibility as in the theorem.

By the theorem, macrostates correspond to pairs (E, S) that can be jointly attained.

The set of macrostates forms the energy-entropy diagram:



Example: The maximum extractable work per copy of a state ρ is given by the horizontal distance to the boundary:



Similarly: analysis of heat engines with finite (but large) reservoirs!

Linear programming

Press release

14 October 1975

The Royal Swedish Academy of Sciences has decided to award the Prize in Economic Science in Memory of Alfred Nobel for 1975 in equal shares to

Professor Leonid Kantorovich, USSR, and Professor Tjalling C. Koopmans, USA,

for their contributions to the theory of optimum allocation of resources.

Optimum Allocation of Resources

Leonid Kantorovich and Tjalling Koopmans have both done their most important scientific work in the field of normative economic theory, *i.e.*, the theory of the optimum allocation of resources. As the starting point of their work in this field, both have studied the problem – fundamental to all economic activity – of how available productive resources can be used to the greatest advantage in the production of goods and services. This field embraces such questions as what goods should be produced, what methods of production should be used and how much of current production should be consumed, and how much reserved to create new resources for future production and consumption.

Linear programming

 $\triangleright~$ Linear programming is appropriate and useful whenever:

- ▷ Resources are arbitrarily divisible.
- ▷ They come in finitely many types.
- ▷ Finitely many basic conversions.
- ▷ In general, all of them fail!
- ▷ So what replaces linear programming?
- $\triangleright\,$ I will explain some results in this direction.

What are resources?

Resources can be **converted** into each other via processes, such as:

timber + nails $\xrightarrow{\text{Carpentry}}$ table

The details vary with the context:

▷ Communication:

noisy channel $\xrightarrow{\text{Channel coding}}$ perfect channel

▷ Thermodynamics:

hot gas + cold gas $\xrightarrow{\text{Carnot process}}$ gas + work

▷ Industrial chemistry:

 $N_2 + 3 H_2 \xrightarrow{Haber process} 2 NH_3$

A mathematical theory of resources

- ▷ Pattern: **convertibility** and **combinability** of resource objects.
- ▷ One investigates questions on catalysis, asymptotic rates,...

Definition (^[4]) An ordered commutative monoid is a structure $(A, +, 0, \ge)$ such that:

$$x + y = y + x \qquad (x + y) + z = x + (y + z) \qquad x + 0 = x$$
$$x \ge y \ge z \quad \Rightarrow \quad x \ge z.$$
$$x \ge y \quad \Rightarrow \quad x + z \ge y + z.$$

 \triangleright No notion of "free resource" or "free operation" is needed.

^[4] Tobias Fritz. "Resource convertibility and ordered commutative monoids". In: Math. Structures Comput. Sci. 27.6 (2017), pp. 850–938.

- ▷ Conceptual insights:
 - \triangleright Catalysts = tools.
 - \triangleright Allowing free use of catalysts \cong having banks to borrow from.
 - $\triangleright~$ Considering asymptotics \cong efficiency of scale and mass production.
 - Asymptotic structure of a resource theory: two convex cones in a vector space.
 - ▷ One for resource objects;
 - ▷ One for resource conversion: $x \ge y$ asymptotically iff x y in this cone.
- ▷ Main open problem: how to **build in approximations?**
- However, we can still exploit the conceptual insights in quantum information theory, for example in quantum thermodynamics!